



# Wireless communications systems security

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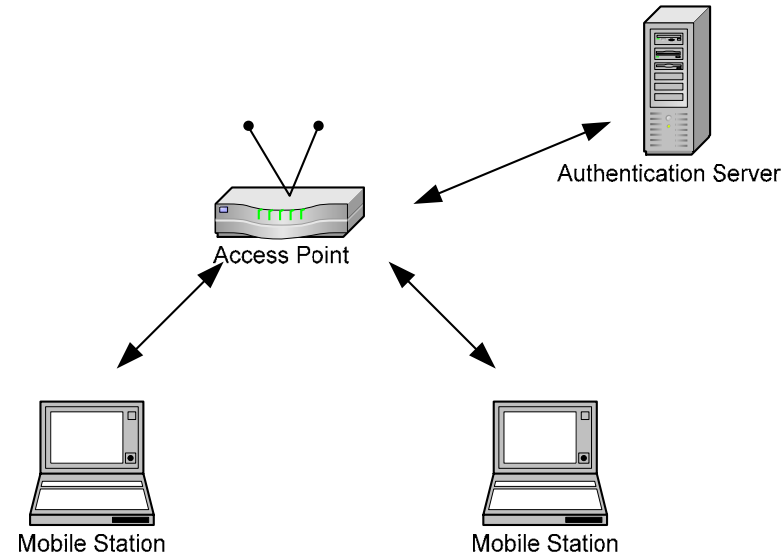


# Agenda

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- Current security solutions in wireless systems (802.11)
- Open problems

# Security Tasks



- Message authentication & privacy
- Node\Base Station mutual Authentication
- Key Management

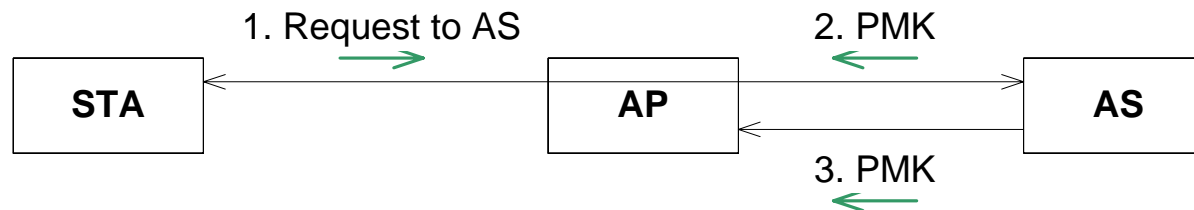
# Message Protection



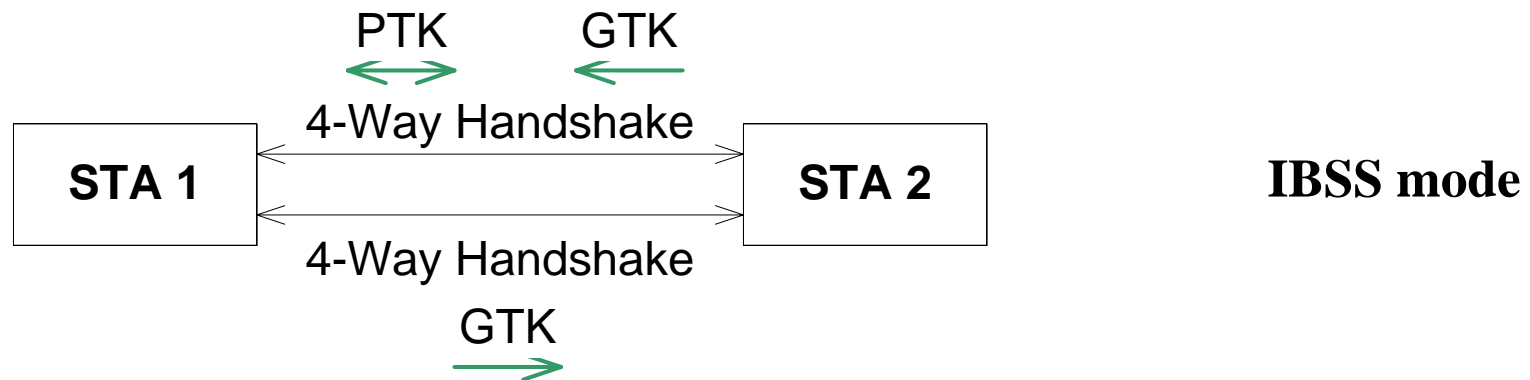
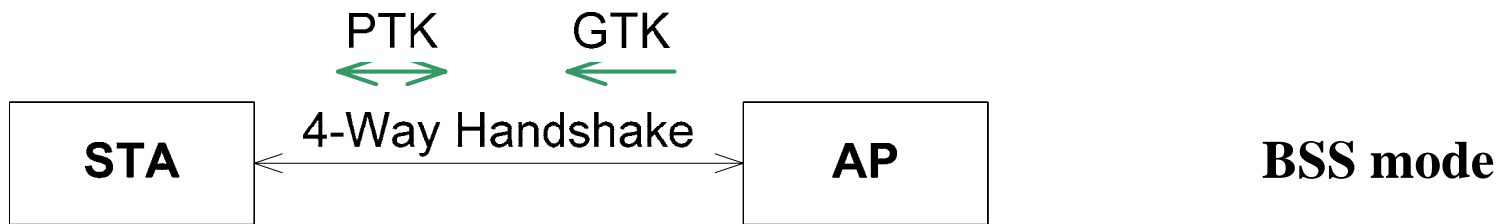
- Message Protection
  - Encryption ( $E^*$ ) (provides Privacy)
  - Packet Number (PN) (provides Freshness)
  - Message Integrity Code (MIC) (provides Authentication)
- Key hierarchy
  - PMK = pair-wise master key
    - PTK = pair-wise transient key (unicast traffic protection)
    - GTK = group temporal key (broadcast traffic protection)

# Key Management

- PMK
  - Administrator provides STA and AP with PSK
    - Only STA and AP are involved
  - Administrator provides STA with key for AS (authentication server)
    - EAP-TLS, EAP-PSK,...
    - Obtaining PMK in a distributed way by using AS (802.1X)



# Obtaining PTK and GTK



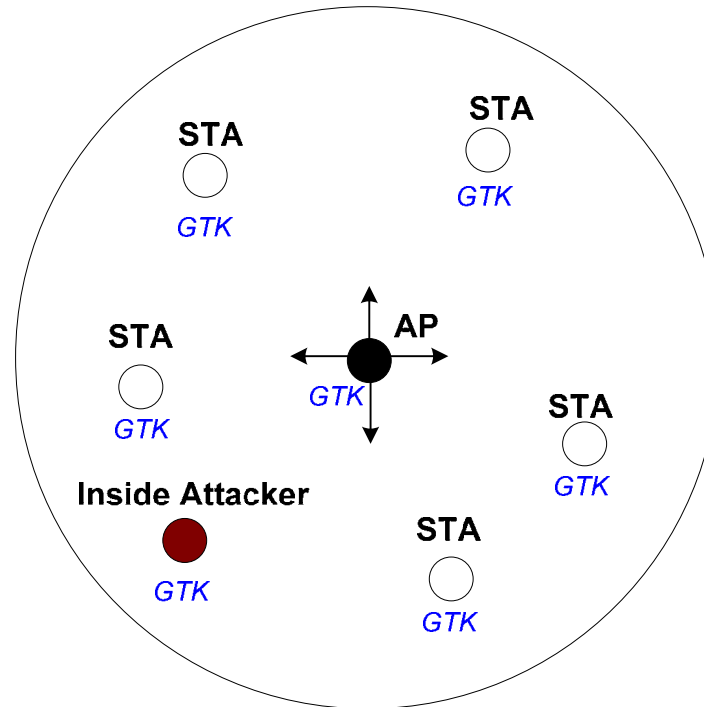


# Open Problems

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- Authentication of broadcast frames
- Authentication of management frames
- Key management in ad-hoc (IBSS) networks

# Broadcast Frames Authentication



- AP distributes GTK to all STAs during association process
- Broadcast data frames are protected by using of GTK
  - Nodes can not identify real source of the broadcast frame
  - Any STA could send broadcast data frames acting as AP





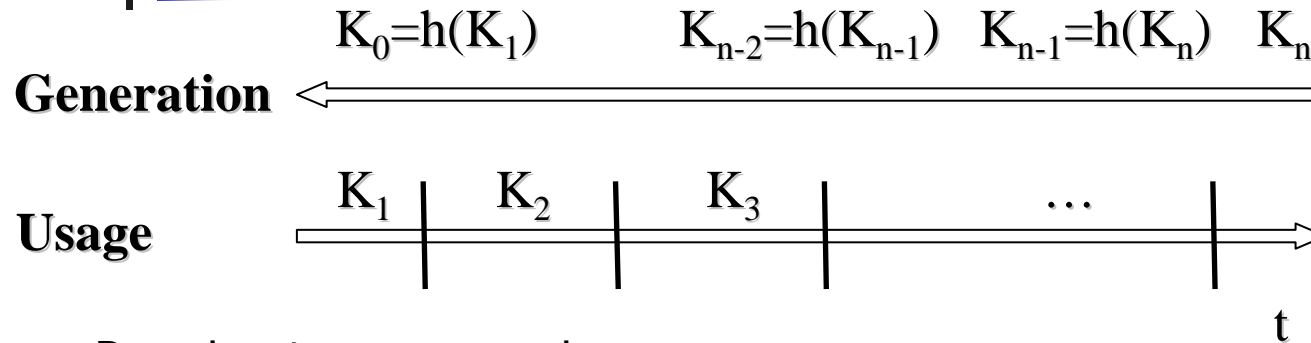
# Possible Solutions

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- Public Key Signature
  - Computational overhead is too big
- MICs on pair-wise keys
  - Network overhead is too big
- One/multi time signature
  - Network overhead is too big
- TESLA
  - Delayed authentication



# TESLA



- $K_{j-1}=H(K_j)$
- $K_0$  - verification key

- Broadcast source sends

$$M \rightarrow M \parallel j \parallel K_{j-1} \parallel \text{MIC}(K_j, j \parallel K_{j-1} \parallel M)$$

for all messages during period  $j$

- Broadcast destinations

- Cache all the frames received during period  $j$
- Verify that  $K_{j-2}=H(K_{j-1})$
- Use  $K_{j-1}$  to validate the MICs of all frames received during the previous period  $j-1$

**Disadvantages:** Delayed authentication, Time synchronization



# Management Frames

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- Mgmt frames are more important for functioning of the network
- Stronger protection should be used
- Broadcast Protection Issues
  - Beacon Protection
  - Deauthentication/deassociation Protection



# Beacon Protection

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- Problem Statement
  - Beacon should be received by STA before authentication/association, i.e. before keys are available
- Possible solution
  - Accept several initial beacons without authentication
  - After authentication/association use TESLA



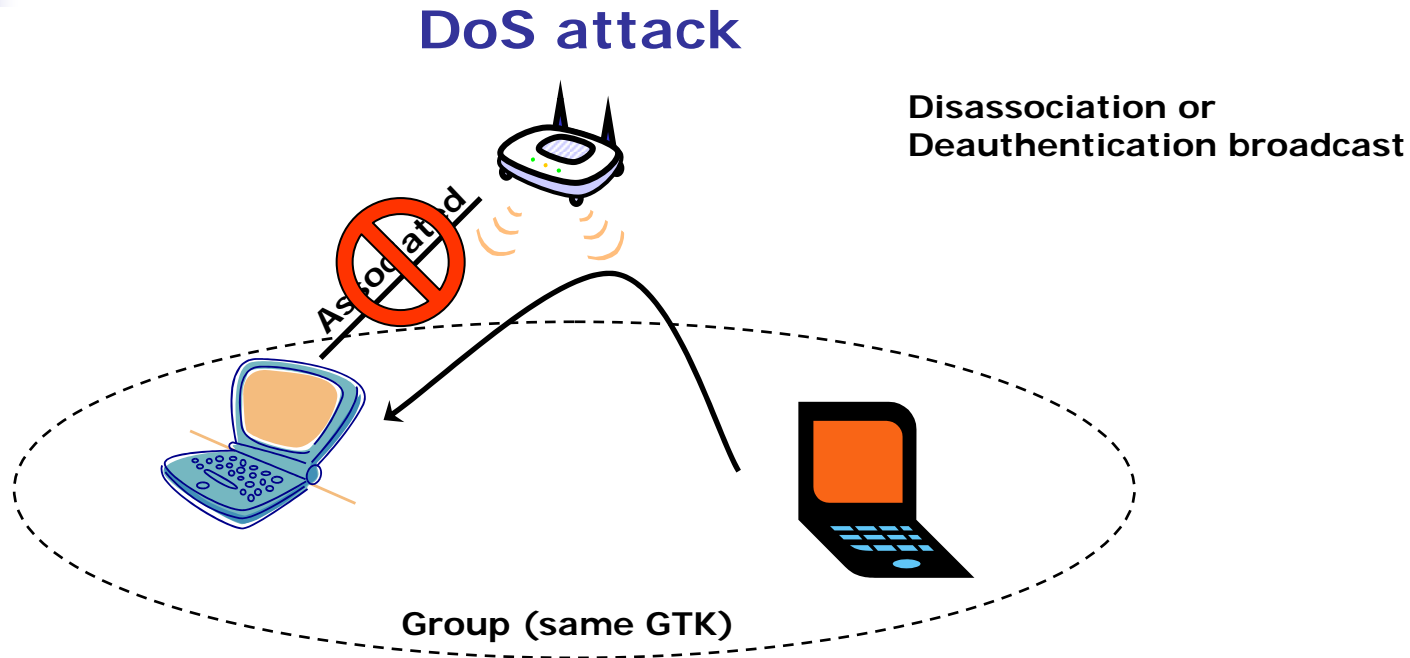
# Beacon Protection

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- Beacon is authenticated using TESLA
- TESLA requires time synchronization
- Time synchronization is provided in beacon

## **Chicken and Egg problem**

# Deauthentication Protection



- Need to protect management frames that a Access Point uses to disassociate or deauthenticate.
- Otherwise attacker can forge such frames which results in a DoS attack.



# Possible Solution

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- AP distributes the commitment value  $CV = \text{hash}(\text{CGTK})$  to STAs
- When AP sends broadcast Disassociation/Deauthenticate, it discloses CGTK
- When an STA receives a protected broadcast, it accepts frame only if  $CV = \text{hash}(\text{CGTK})$
- This broadcast frame was successfully protected, because it is sent only once per session.
  - One-time signature



# Public Key Infrastructure (PKI)

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- Binding between PK and user
  - User certificate (CERT)  
CERT = <Name, Tstart, Tend, PK, CA Signature>
- Certificate management
  - Issuing, renewal, revocation
  - Done by trusted third party Certificate Authority (CA)





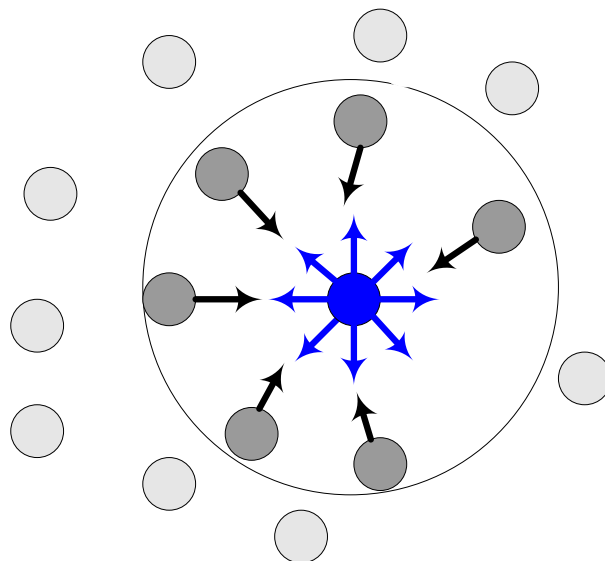
# Key Management

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- Centralized certification authority (CA), which provides STAs with certificates, exists in scenarios discussed previously
- There are scenarios (e.g. MANET), where such authority is not available (no fixed infrastructure)
  - To use certificate in self-organized networks we need to distribute the functionality of CA
    - Distribute signing procedure

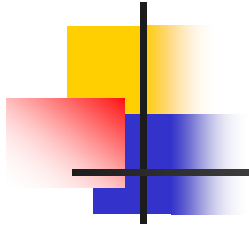
# Distributed Certification Authority

- Distributed signing procedure => distributed secret key
- Secret Sharing  $(t,n)$ -scheme
  - $SK_{CA}$  is shared among all  $n$  nodes
  - $t < n$  nodes can make calculation with  $SK_{CA}$
  - Less than  $t$  nodes can not



## Threshold (7,17) - scheme

- Coalition member
- Requester



**Thank You**